

Name:

NetEcon final exam

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For each question, check all boxes corresponding to correct answers. There may be zero, one or several.

Advice: Read the questions carefully!

1. Consider a game with n users sharing the same communication link. Each user i chooses a non-negative consumption x_i and receives a utility $u_i(x_i, x_{-i}) = f(x_i) - (x_1 + \dots + x_n) - p_i$ where $f(\cdot)$ is an increasing concave function and p_i is the price for user i .

- ☐ The social welfare is $\sum_i f(x_i)$.
- ☐ If $p_i=0$ for each user i , the price of Anarchy is one.
- ☐ The price of anarchy is always larger or equal to one.
- ☐ If $p_i=nx_i$, for all user i , at NE each user chooses a socially optimal level of consumption.

2. Consider the following two-players game:

		P2	
		A	B
P1	A	5, 2	-1, -1
	B	-1, -1	2, 5

- ☐ It is a potential game with potential f such that $f(a, a)=5$, $f(a, b)=2$, $f(b, b)=5$, $f(b, a)=-1$.
 - ☐ There exists an infinity of functions f such that it is a potential game with potential f .
 - ☐ Every finite potential game has a pure strategy Nash equilibrium.
 - ☐ Every finite potential game has a fully mixed Nash equilibrium.
3. Consider a 2-players attacker defender game. The attacker has 2 actions, attack (a) or not-attack (na) and the defender has 2 actions, monitor (m) or not monitor (nm). The payoffs are (with $\alpha_c > 0$, $\alpha_f > 0$, $\alpha_s > 0$, $\beta_c > 0$, $\beta_s > 0$):

		defender	
		m	nm
attacker	a	$-\beta_c, \alpha_c$	$\beta_s, -\alpha_s$
	na	$0, -\alpha_f$	$0, 0$

- ☐ Every finite game has a mixed strategy Nash equilibrium.
 - ☐ The maxmin strategy (or safe strategy) and the Nash equilibrium always coincide.
 - ☐ At the Nash equilibrium of this game, the attacker is choosing “a” with probability $1/(\alpha_c + \alpha_f + \alpha_s)$.
 - ☐ At the Nash equilibrium of this game, the attacker’s expected payoff is strictly positive.
4. Auctions. We consider auctions with a single item, where there is one seller and there are $n > 1$ buyers with independent identically distributed private value.
- ☐ A second-price auction is equivalent to an open descending auction.
 - ☐ Running an auction is always the optimal way of selling the good, regardless of the information of the seller about the buyers’ values.
 - ☐ If a seller who wants to maximize his expected revenue through a second-price auction has value zero for the good, he should set a reserve price of zero.
 - ☐ In a first-price auction, at the symmetric increasing Nash equilibrium, every user bids less than his true value.